## Spring Semester 2005 Topics in Computation Theory (CS700) Discrete Geometry Homework 5

This homework is due on *Wednesday* May 4, at the beginning of the (extra) class at 4:00 p.m.

On the top of the first sheet that you turn in, please put (a) your name and student number, (b) how much time you spent working on the homework, and (c) a little table with your self-evaluation as explained on the course webpage.

- 1. Prove that the region reg(p) of a point p in the Voronoi diagram of a finite point set  $P \subset \mathbb{R}^d$  is unbounded if and only if p lies on the surface of conv(P).
- 2. (a) Let C be any circle in the plane x<sub>3</sub> = 0. Show that there exists a half-space h such that C is the vertical projection of the set h ∩ U onto the plane x<sub>3</sub> = 0, where U = {x ∈ ℝ<sup>3</sup> | x<sub>3</sub> = x<sub>1</sub><sup>2</sup> + x<sub>2</sub><sup>2</sup>} is the unit paraboloid.
  (b) Consider n arbitrary circular disks K<sub>1</sub>,..., K<sub>n</sub> in the plane. Show that the

boundary of the union  $\bigcup_{i=1}^{n} K_i$  consists of O(n) circular arcs.

- 3. Define a "spherical polytope" as an intersection of n balls in  $\mathbb{R}^3$  (such an object has facets, edges, and vertices similar to an ordinary convex polytope).
  - (a) Show that any such spherical polytope in  $\mathbb{R}^3$  has  $O(n^2)$  faces. You may assume that the spheres are in general position.
  - (b) Show that the intersection of n unit balls (that is, balls of radius one) has O(n) complexity.
- 4. Prove that if  $X \subset \mathbb{R}^d$  is a set where every two points have distance 1, then  $|X| \leq d+1$ .
- 5. Give an algorithm to solve Sylvester's problem: given a set of n points in the plane, the algorithm should output a line containing exactly two points (or report that all points lie on a line). Can you achieve sub-quadratic running time? A "perfect" solution should have running time  $O(n \log n)$ . Bonus points will be awarded for a linear-time algorithm!