k-COLORING:

Instance: A graph G.

Question: Can the vertices of G be colored with k colors such

k-Coloring $\leq (k+1)$ -Coloring

that no two vertices of the same color are adjacent?

(Note that k is part of the problem, not an input.)

Lemma: 2-COLORING can be solved in linear time.

Graph Coloring

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Independent Set and Coloring

3-Coloring \leq IndependentSet



n vertices

3n vertices, k = n

(v, u, i)

3-Coloring \Leftrightarrow independent set of size n





There are three possible ways for a Hamiltonian cycle to visit these four vertices.

(u, v, i)

Vertex Cover and Hamiltonian Cycle

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Vertex Cover and Hamiltonian Cycle

 $\mathrm{VertexCover} \leq \mathrm{HamiltonianCycle}$



There are three possible ways for a Hamiltonian cycle to visit these four vertices.



 $VERTEXCOVER \leq HAMILTONIANCYCLE$



There are three possible ways for a Hamiltonian cycle to visit these four vertices.

G has a vertex cover of size k if and only if G' has a Hamiltonian path.

Let u_1, u_2, \ldots, u_k be the vertices of the vertex cover C. Then the Hamiltonian path starts in cover vertex 1, visits the vertex chain of u_1 , goes to cover vertex 2, visits the vertex chain of u_2 , and so on, until returning to cover vertex 1.



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Example Cover







Subset Sum

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Summary

SUBSETSUM:

Instance: A set X of positive integers and an integer t. Question: Does X have a subset whose elements sum to t?

 $VERTEXCOVER \leq SUBSETSUM$

Number edges from 0 to m-1.

Our set X contains $b_i = 4^i$ for each edge i, and a_v for each vertex v:

$$a_v = 4^m + \sum_{i \in \Delta(v)} 4^i.$$

The target sum $t \mbox{ is }$

$$t = k \cdot 4^m + \sum_{i=0}^{m-1} 2 \cdot 4^i.$$

3-Coloring \leq IndependentSet

3-Coloring \leq Planar3Coloring

INDEPENDENTSET \leq VERTEXCOVER

INDEPENDENTSET \leq CLIQUE

 $VERTEXCOVER \leq SETCOVER$

 $VertexCover \leq HamiltonianCycle$

 $VERTEXCOVER \leq SUBSETSUM$

SubsetSum \leq Partition

HamiltonianCycle \leq HamiltonianPath

 ${\rm HamiltonianPath} \leq {\rm LongestPath}$