## KAIST CS206

Sorting

Sorting problem: Given a list a with n elements possessing a total order, return a list with the same elements in non-decreasing order.

The sorting problem is perhaps the most fundamental problem in algorithms.

We can sort any kind of element that can be compared (int, float, str). In other words, we require a total order on the elements.

There are many direct applications of sorting (catalogs, reports, file listings, etc.)

## KAIST CS206

Sorting as a tool

duplicates2.pv

There are also many indirect applications of sorting. For instance, algorithms can often be made faster by first sorting the data.

```
def has_duplicates_sorted(a):
  for i in range(len(a)-1):
    if a[i] == a[i+1]:
       return True
  return False
```

```
def has_duplicates(a):
    return has_duplicates_sorted(sorted(a))
```

Sorting + linear time!

This implementation works, but it creates a lot of lists, copies

a lot of data, and could cause a runtime stack overflow...

KAIST CS206	Selection sort	KAIST CS206	Selection sort
It's easy to find the minimum of $n$ numbers:		<pre>def selection_sort(a):     if len(a) &lt;= 1:         return a     k = find_min_index(a)     b = selection_sort(a[:k] + a[k+1:])     return [a[k]]+b</pre>	selection0.py
<pre>def find_min_index(a):</pre>			
mindex = 0			
for k in range(1, len(a)):			
<pre>if a[k] &lt; a[mindex]:     mindex = k</pre>			
return mindex		What is the running time?	

This gives immediately a sorting algorithm:

- A list with zero or one element is already sorted.
- Otherwise, find the minimum element, and recursively sort the remaining n-1 elements.
- Concatenate the minimum and the sorted remaining elements.

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In-place Sorting

Sorting problem: Given a list a with n elements possessing a total order, return a list with the same elements in non-decreasing order.

Often we no longer need the original, unsorted data.

In-place Sorting: Given a list a with n elements possessing a total order, rearrange the elements inside the list into non-decreasing order.

Saves a lot of memory for huge data. Ideally we want to do this without creating any other list.

In Python:

- sorted(a) returns a sorted copy of a.
- a.sort() sorts the list a in-place.

# **KAIST** CS206 In-place selection sort with iteration

```
def find_min_index(a, i): Find index of minimum in a[i:]
mindex = i
for k in range(i+1, len(a)):
    if a[k] < a[mindex]:
    mindex = k
return mindex

def selection_sort(a):
    n = len(a)
for i in range(0, n-1):
    k = find_min_index(a, i)
    t = a[i]
    a[i] = a[k]
    a[k] = t</pre>
```

This uses only one list (in-place) and cannot have stack overflow, but the running time is still  ${\cal O}(n^2).$ 

## KAIST CS206

#### In-place Selection sort

def find\_min\_index(a, i): Find index of minimum in a[i:]
 mindex = i
 for k in range(i+1, len(a)):
 if a[k] < a[mindex]:
 mindex = k
 return mindex</pre>

## KAIST CS206

Insertion sort

insertion0.pv

Let's do it the other way round: Sort n-1 elements first, then insert the last element into the sorted sequence.

```
def insertion_sort(a):
    if len(a) <= 1:
        return a
    b = insertion_sort(a[:-1])
    k = sorted_linear_search(b, a[-1])
    b.insert(k, a[-1])
    return b</pre>
```

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#### In-place insertion sort

insertion1.pv

This is not tail-recursion, but we can still easily make it iterative.

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Iterative in-place insertion sort

correctness of the program.

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Bubble Sort

bubble1.py

Similar to selection sort, we bring the largest element to the end:

```
def bubble_sort(a):
  for last in range(len(a), 1, -1):
    # bubble max in a[:last] to a[last-1]
    for j in range(last-1):
        if a[j] > a[j+1]:
        t = a[j]
        a[j] = a[j+1]
        Bubble-up phase
        a[j+1] = t
```

If nothing happens during a bubble-up phase, we are done!

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#### Bubble sort with early termination

```
We stop when nothing happens in one phase.
def bubble_sort(a):
  for last in range(len(a), 1, -1):
    # bubble max in a[:last] to a[last-1]
    flipped = False
    for j in range(last-1):
      if a[j] > a[j+1]:
                            Effective if the list is already
        flipped = True
                            (nearly) sorted.
        t = a[j]
        a[j] = a[j+1]
        a[j+1] = t
                            What is the worst case
    if not flipped:
                            running time?
      return
```